

Binary Search Problem (D&C Approach)

- The algorithm is used to look up a word in a dictionary or a name in a telephone directory.
- Let $T [1..n]$ be an array sorted in ascending order.
- Let x be some item. The problem consists of finding x in the array. If x is not in the array, then find the position where it might be inserted.
- **1st Approach: Sequentially**
 - Function sequential ($T [1..n], x$)
{Sequential search for x in the array T }
for $i \leftarrow 1$ to n do
 if $T[i] \geq x$ then return i
return $n + 1$
- Worst case $\Omega(n)$ and Best case $O(1)$, time is $\Theta(n)$

Binary Search Problem (2nd Approach)

- Use D & C Approach:
 - To Speed Up the search, look for x either in first half or in the second half of an array.
 - Let $k = \lfloor n/2 \rfloor$; If $x \leq T[k]$, then search x in $T[1 \dots k]$; otherwise search in $T[k+1 \dots n]$

Binary Search Problem (2nd Approach)

- **Algorithm:**

Function binsearch (T[1...n], x)

if $n=0$ or $x > T[n]$ then return $n+1$
else return binrec(T[1...n], x)

Function binrec(T[i...j], x)

{ Binary search for x in subarray T[i...j] with promise that
 $T[i-1] < x \leq T[j]$ }

if $i=j$ then return i

$k \leftarrow (i+j)/2$

if $x \leq T[k]$ then return binrec(T[i...k], x)

else return binrec(T[k+1...j], x)

Binary Search Problem (2nd Approach)

X=12

1	2	3	4	5	6	7	8	9	10	11	
-5	-2	0	3	8	8	9	12	12	26	31	X ≤ T[k]
i					k					j	No
						i		k		j	Yes
						i	k	j			Yes
						ik	j				No
							ij				i=j ; stop

Worst & Average Case: O (log N), Best Case: O(1)